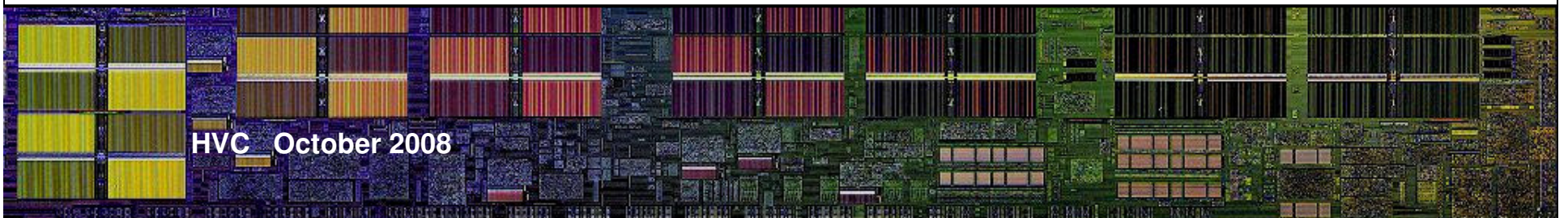


Moore's Law v. Verification Complexity

Jason Baumgartner

Thanks for contributions: Wolfgang Roesner,
Ralf Fischer

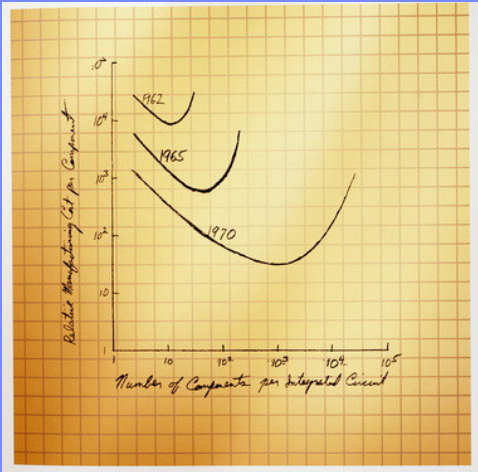
IBM Corporation



Outline

- What is Moore's Law?
- Verification Complexity
- Does Moore's Law hurt Verification Complexity?
 - No
 - Yes
- Can we cope with this complexity?
 - Yes
 - No
- Open problems: *despair!!! vs hope???*

What is Moore's Law?



transistors per IC for minimum cost has increased at roughly a factor of two per year



there is no reason to believe it will not remain nearly constant for at least 10 years

G. Moore, "Cramming More Components Onto Integrated Circuits," Electronics Magazine 1965

What *Isn't* Moore's Law?

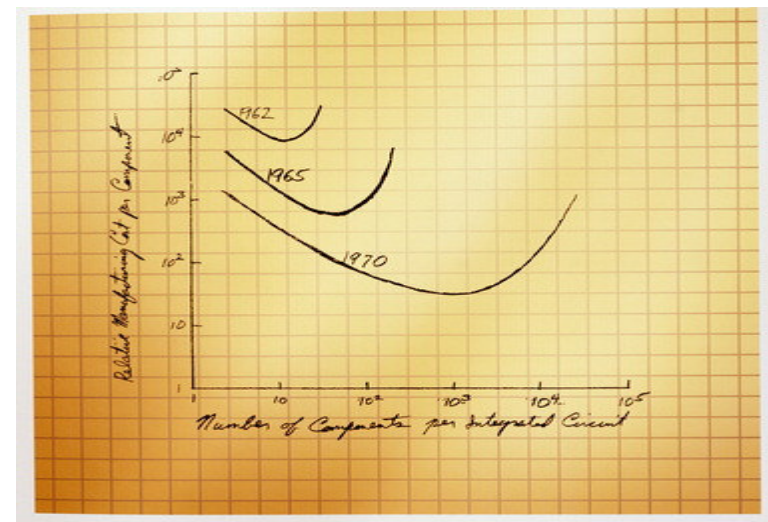
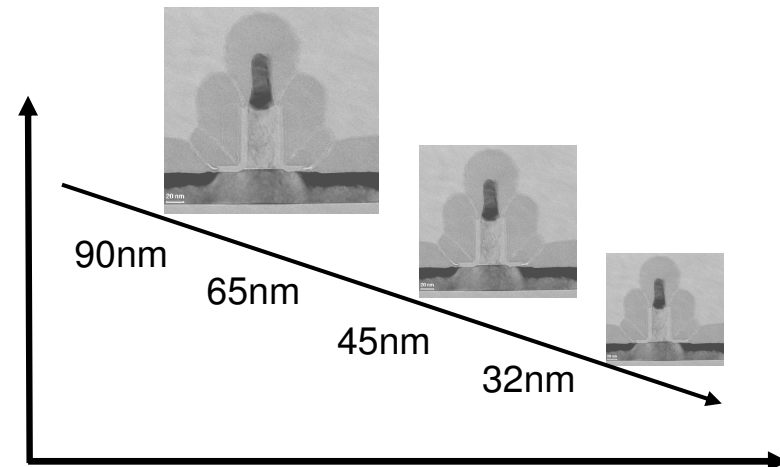
- Attributed to virtually all exponentially-growing computing metrics
 - Circuit speed
 - Computing power (MIPS, GFlops, ...)
 - Storage capacity
 - Network capacity
 - Pixel density
 - ...
- Strictly speaking, these are *not* part of Moore's original observation
 - IC complexity for *minimum cost*

What *Isn't* Moore's Law?

- Nonetheless, all refer abstractly to the same trend
 - Device miniaturization
 - Reliable integration “growth”
 - “Circuit and device cleverness”
- We deliberately abuse notation w.r.t. “*Moore's Law*” in this talk

Design Characteristics of “Moore’s Law”

- Smaller: *miniaturization*
 - Devices and transistors
- Cheaper
 - Per *transistor*
 - Not necessarily per *product*
- Faster
 - If software developers are willing :-)



Design Characteristics of Moore's Law

- Bigger!

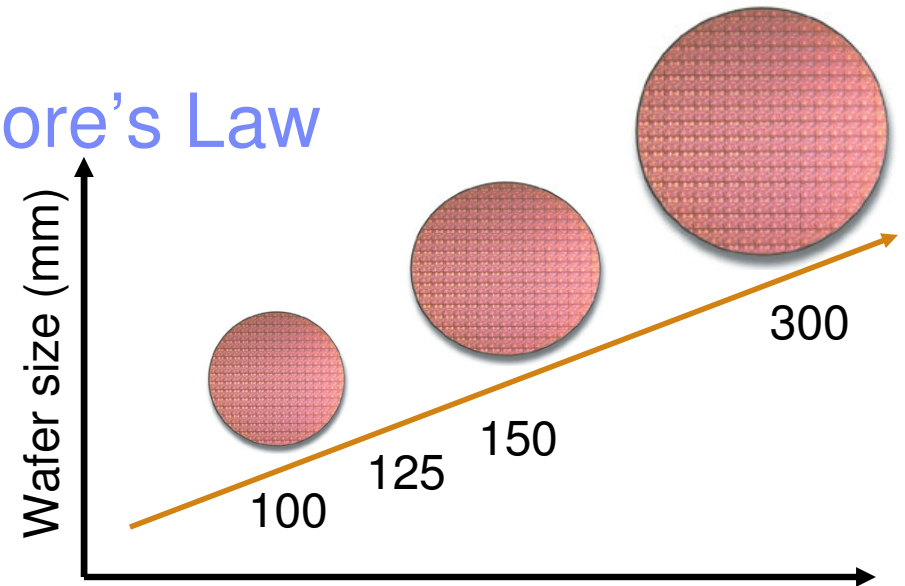
- Chips, wafers, ICs, networked systems, ...

- More complex!!

- More “smaller” devices crammed onto chip / IC
- Scale-up of functionality: datawidth, memory size, ...

- Hotter!!!

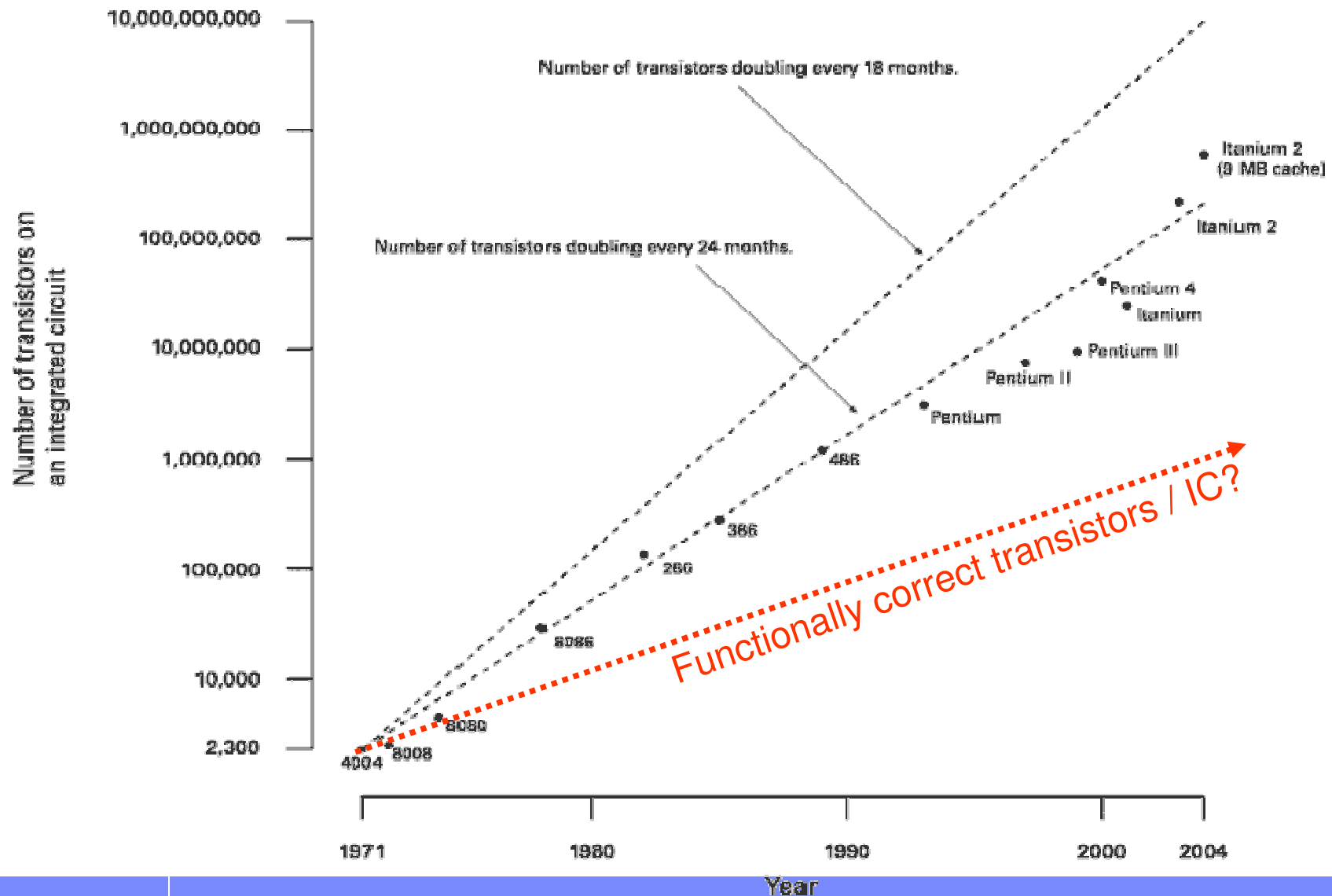
- Since bigger & faster!



Design Characteristics of Moore's Law

- Comparable functionality in smaller / cheaper package?
- No! Cram *more* into a bigger package
- Harder to verify!!! ??? !?*\$%#@!!
 - Thankfully, “device complexity” cannot *afford* to be “as great as possible”

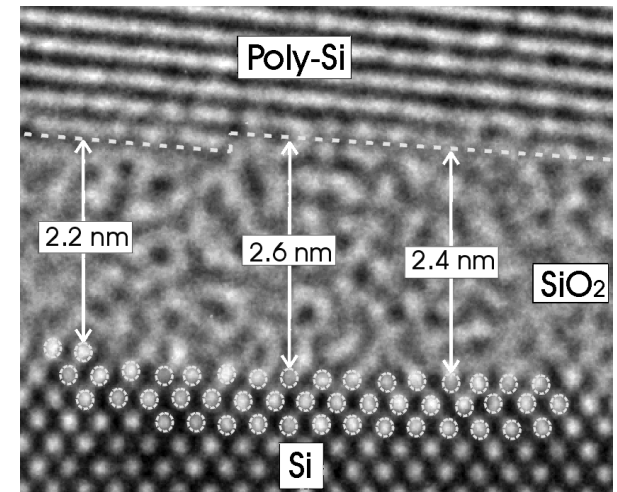
Longevity: for at least 10 years, indeed!



The End is Near! (Is it?)

- *Moore himself* was one of his harshest critics

- Disappearing “circuit cleverness” 1975
- Lack of demand for VLSI 1979
 - *The Death of Cost Effectiveness*



- *No exponential is forever* - must hit limitations of physics?

- Can we miniaturize (and design w.r.t.) quantum particles? Hmmm...

- Note trends on massive parallelization (e.g. BlueGene), 3D chips, biological computing, quantum computing, ...

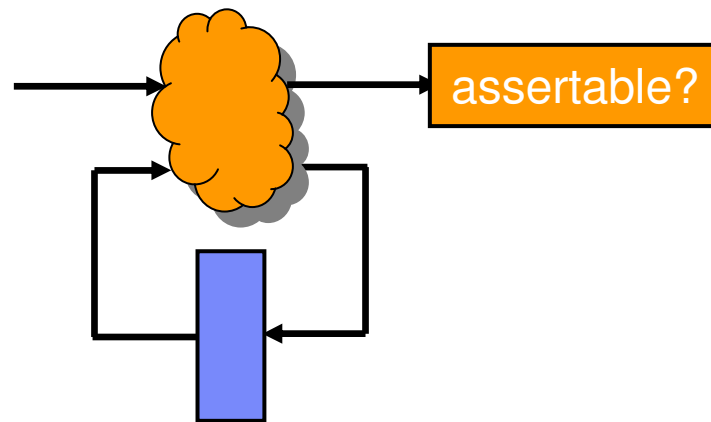
- Who knows?
- Will global warming (or out-of-control particle accelerator) finitize “forever”?

What is “Hardware Verification”?

- Chips often designed in Hardware Description Language (HDL)
- HDL taken through compile, physical design steps to fabrication
- Many facets to verifying a chip
- *Logic verification* is the primary focus of this talk
- *Does the HDL of the chip produce correct computations?*
 - E.g., the FPU generate IEEE-compliant results?
- Will later touch upon correctness of post-HDL flow

What is “Hardware Verification”?

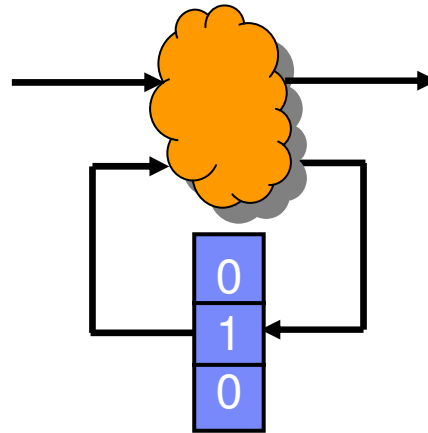
- May often represent verif problem using a *sequential netlist*
 - Correctness properties may be synthesized into simple assertion checks



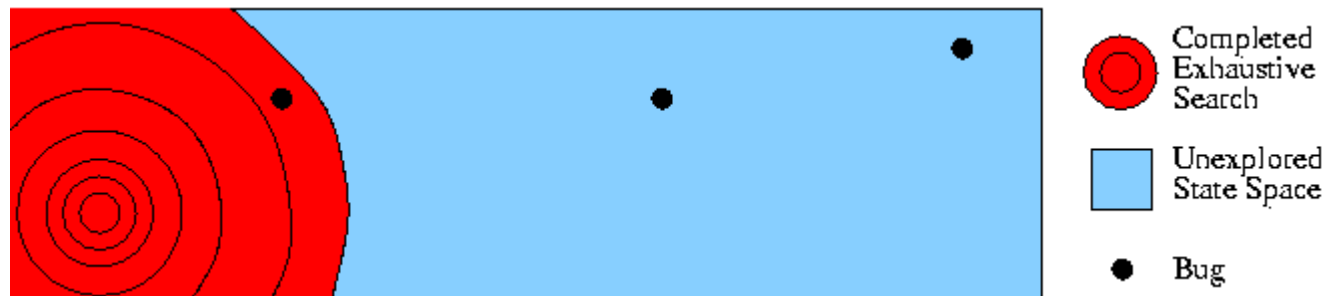
- Netlist composed of primary inputs, combinational gates, sequential elements (latches, RAM, ...)
- *Sequential* refers to ability to “remember” past computations

Verification Complexity

- A *state* is a valuation to the sequential elements of the design

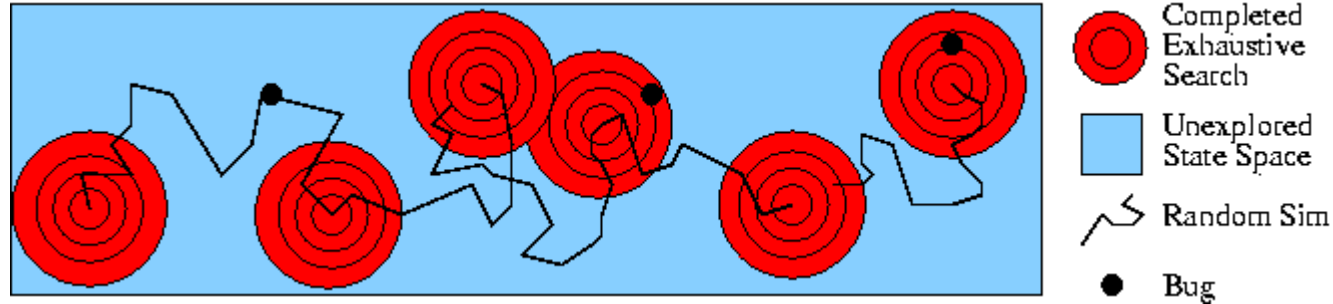


- Exhaustive (formal) verification generally requires analysis of *reachable* states



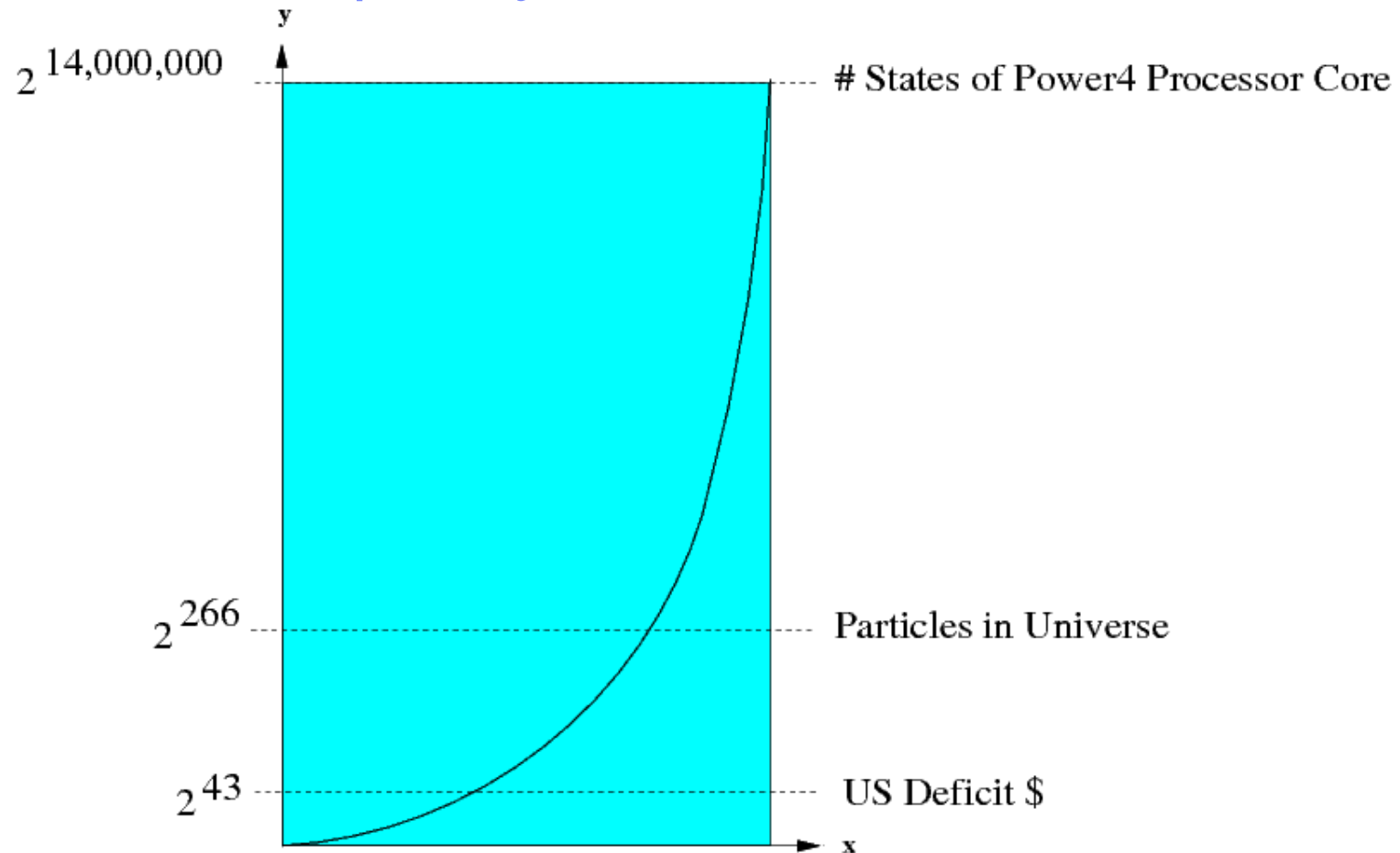
Verification Complexity

- Formal verification generally requires analysis of *reachable states*
 - Falsification may only require exploring a *subset*



- Some proof techniques leverage fast analysis of a *superset*
 - Induction: can design transition from a good state to a bad state?

Verification Complexity



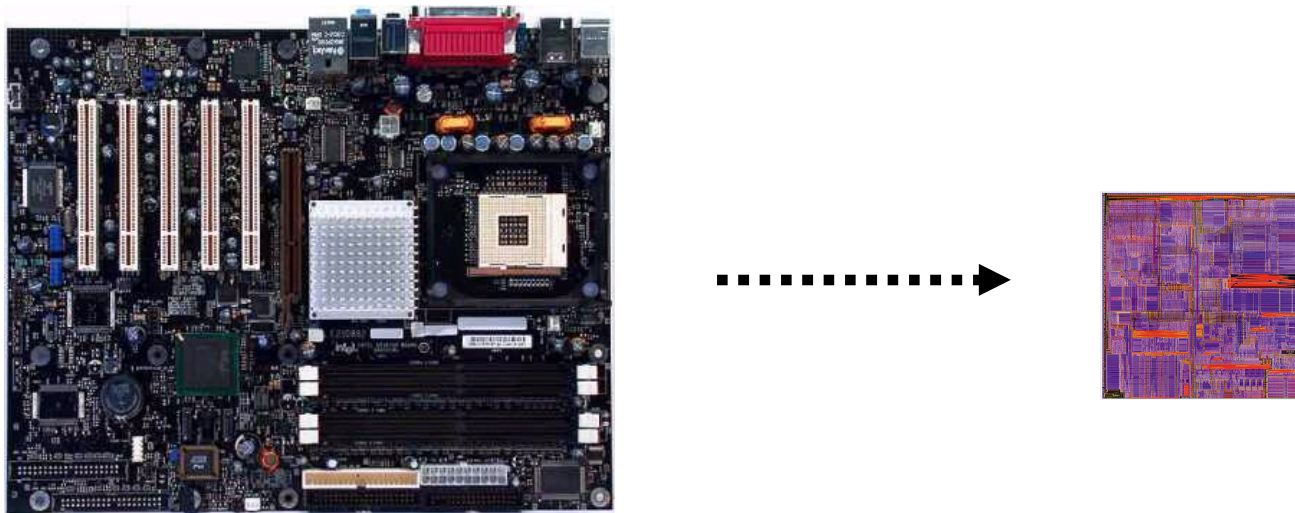
- *Explicit* state enumeration is intractable
 - *Symbolic* analysis often superior, though still capacity-gated

Moore's Law v. Verification Complexity

- # Components per IC doubles every ~2 years
- Verification thus *appears* to grow exponentially more complex
 - Compounded by use of *today's* computers to verify *tomorrow's* designs
- Is this *necessarily* the case?
- Let us revisit how this capacity tends to be used
 - Moore's *Heirlooms*

Moore's Heirlooms: *Integration*

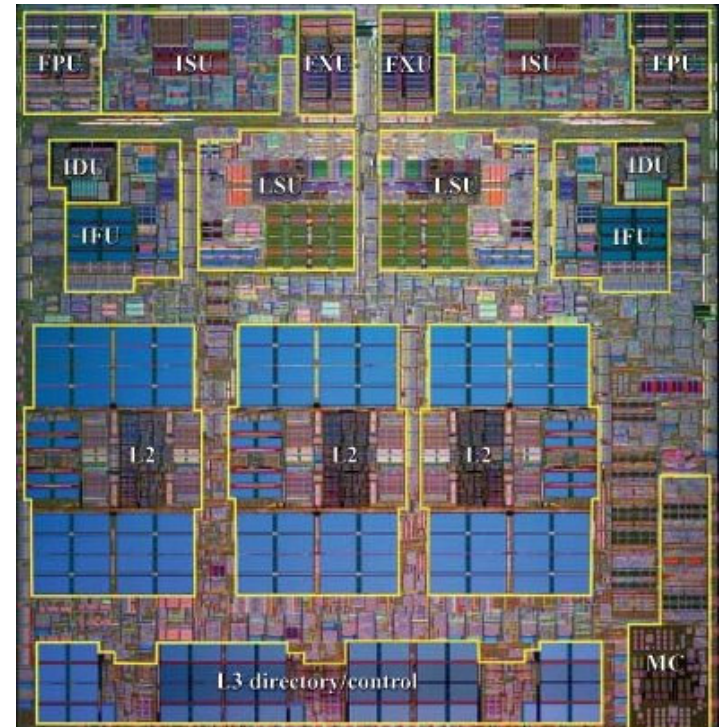
- Integration of more devices on chip
 - System on a Chip: more components+functionality moved on-chip
 - Caches are moving on-chip



- Lowers packaging costs and power, increases speed
- “Moving” components: no negative impact to verif complexity

Moore's Heirlooms: *Modularity*

- Additional execution units
 - Multiple FPUs, FXUs, LSUs, ...
- Additional cores
 - POWER4 is 2 core; POWER7 is 8 core
- No additional *component* verif complexity
- Overall *system* complexity may increase
 - Hardware, software, or both
 - More concurrency, # interfaces
 - Some aspects may be covered by higher-level verification



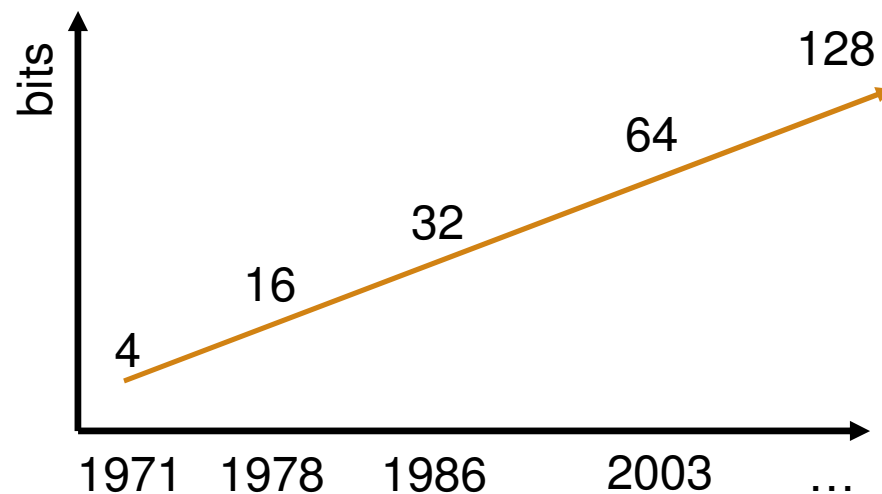
Moore's Heirlooms: *Specialized Hardware*

- SW function moves to hardware
 - Vector units, encryption
- *Diversified* modularity
 - Cell processor: 8 Synergistic Processing Elements in addition to a Power processor
- May not increase verif complexity
 - “Only” *more* components to verify
- Though nonetheless difficult to verify!
 - Move verification burden from SW to HW?



Moore's Heirlooms: *Increased Operand / Data Width*

- Operand width has grown substantially
 - Mainstream (vs *mainframe*!) processors



- Many processors have emulated 128-bit data support for decades
 - SW + specialized HW atomically manages narrower computations

Moore's Heirlooms: *Increased Operand / Data Width*

- Does increased data width increase verification complexity?
 - Sometimes “no” !!!
- Data routing checks are not necessarily more complex
 - Some checks may be **bit-sliced**; *linear* verification scaling
 - **Word / vector** reasoning techniques scale well *when applicable*
 - UCLID, SMT, uninterpreted functions
 - Verification **reduction techniques** have been proposed to automatically shrink widths to facilitate a broader set of algorithms
 - Control / token nets, Bjesse CAV'08

Moore's Heirlooms: *Increased Operand / Data Width*

- Does increased data width increase verification complexity?
 - Sometimes “yes” !!!
- What about correctness of *computations* on the operands?
 - *Optimized* arithmetic / logical computations are not simple $+$ $=$ $*$ $/$ $<$ $>$
- Consider IEEE Floating Point Spec

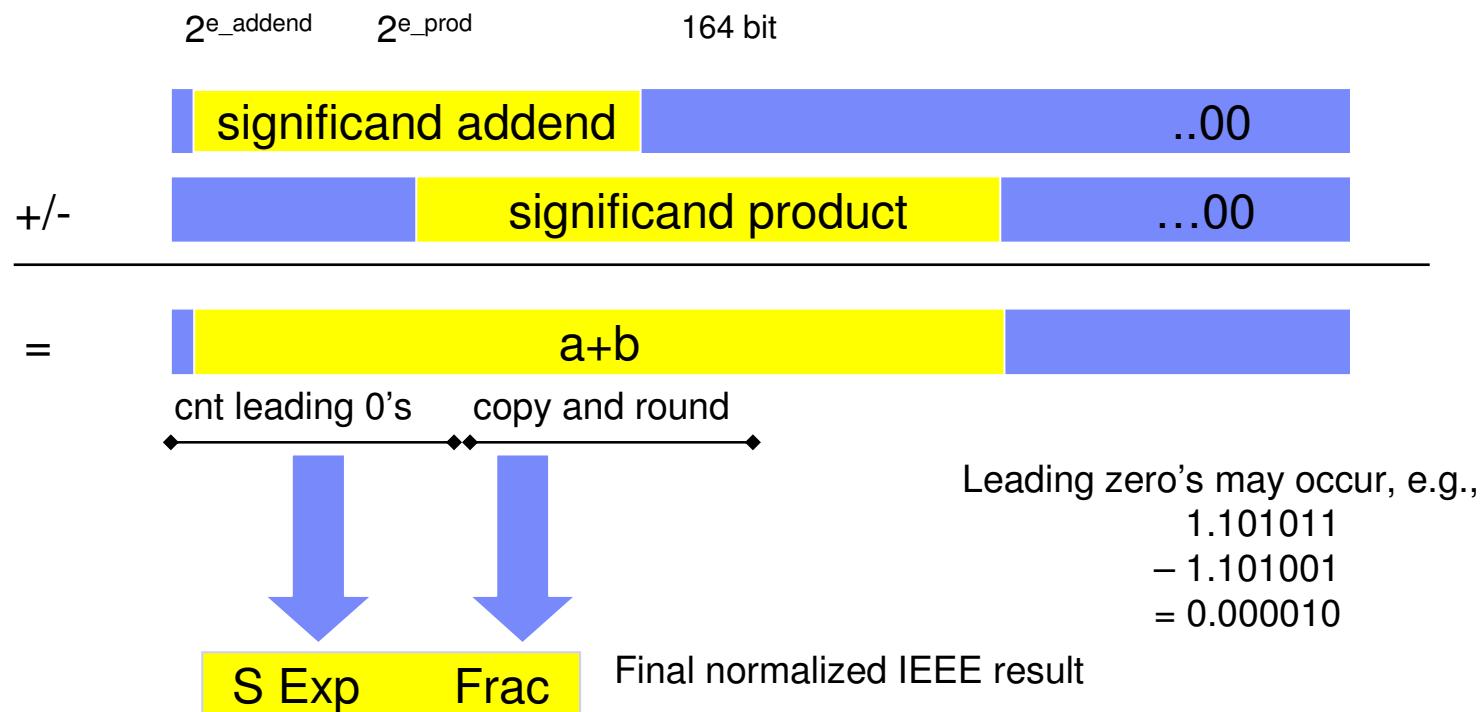
	Single Precision	Double Precision	Quadruple Precision
Width	32	64	128
Exponent bits	8	11	15
Significand bits	23	52	112

Floating-Point Verification

- Floating point number format: $S * B^E$
 - S: *Significand*, e.g. 3.14159
 - B: *Base*, here $B=2$
 - E: *Exponent*, represented relative to predefined *bias*
 - Actual exponent value = $bias + E$
- A *normalized* FP number has Mantissa of form 1.?????
- Aside from *zero* representation
- *Fused multiply-add* op: $A*B + C$ for floating point numbers A,B,C
 - C referred to as *addend*
 - $A*B$ referred to as *product*

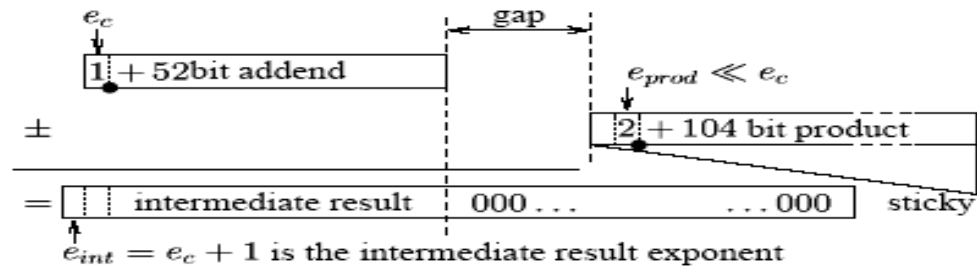
Floating-Point Verification: Double Precision

- Optimized FPU: 15 000 lines HDL
- IEEE-compliant reference model for double precision: 500 lines HDL

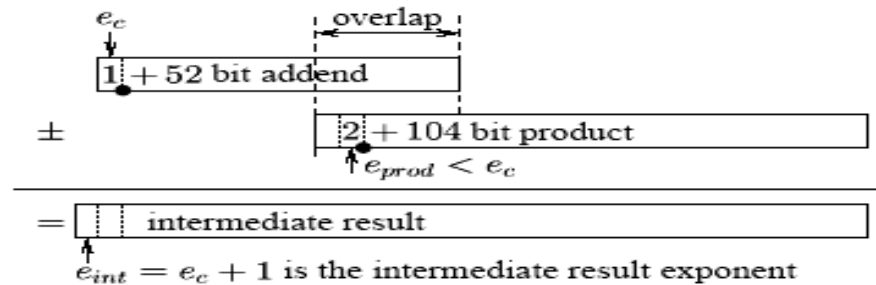


Floating-Point Verification: Double Precision

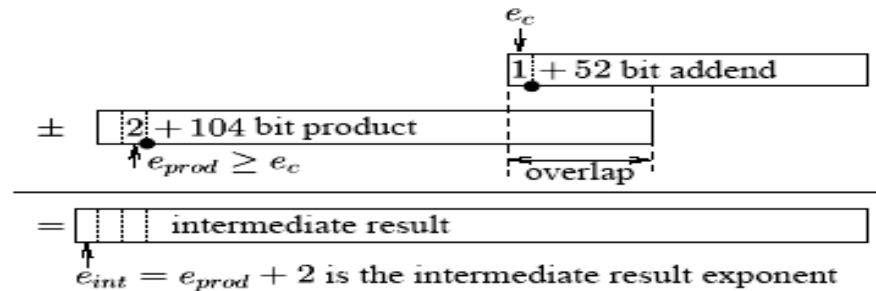
- Direct equiv check between reference, implementation is computationally intractable
 - Use *case-splitting* strategy!
- Four distinct categories of case splits used
 - Based on difference between product, addend exponent



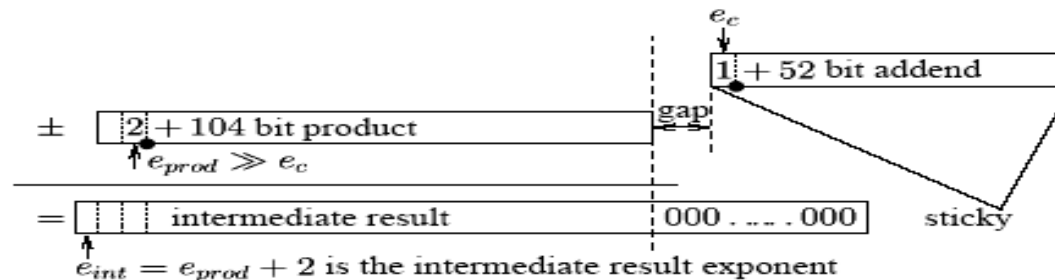
(a) Farout left: the addend is much larger than the product; there is no overlap, the product becomes a sticky bit.



(b) Overlap left: the addend is larger than the product; the product overlaps with the right part of the addend.



(c) Overlap right: the product is larger than the addend; the addend overlaps with the right part of the product.



(d) Farout right: the product is much larger than the addend; the addend becomes a sticky bit.

Floating-Point Verification: Double Precision

- Case split to fix product exponent $C_\delta := (e_a + e_b - bias = e_c + \delta)$

$\delta = e_{prod} - e_c$ where $e_{prod} (= e_a + e_b - bias)$ is the product exponent
and e_c is the addend exponent

- Normalization shifter is used to yield a *normal* result

- Depends upon # number of leading zeros of intermediate result

$C_{sha} := (sha = X)$ for all 106 possible shift amounts;

$C_{sha / rest} := (sha > 106)$ to cover the remaining cases (trivially discharged)

- Define a secondary case-split on normalization shift

- Constraint defined directly on shift-amount signal (sha) of Ref-FPU
- *Sha* is 7-bit signal (double-precision) to cover all possible shift amounts

Floating-Point Verif: Double to Quad Precision

- Double precision
 - ~585 total cases to check
 - Each tractable using BDDs
- Quad precision
 - ~1244 total cases
 - *None* are practical using BDDs
- Denormal operands require additional case split on input normalization shifter
 - # Cases increases three orders of magnitudes double-to-quad

Moore's Heirlooms: *Increased Operand / Data Width*

- Error Code Detection / Correction (ECC) logic becomes substantially more complex w.r.t. data width
 - Byproduct of transistor miniaturization: *soft errors!*
 - Increasingly mandate ECC logic
 - Along with increasingly elaborate ECC algos to handle *more* error bits
- Emerging encryption HW similarly explodes in complexity w.r.t. data width
- Overall: *Does increased data width increase verif complexity?*
 - Sometimes dramatically !?*\$%*#@!!

Moore's Heirlooms: *Increased RAM Depth*

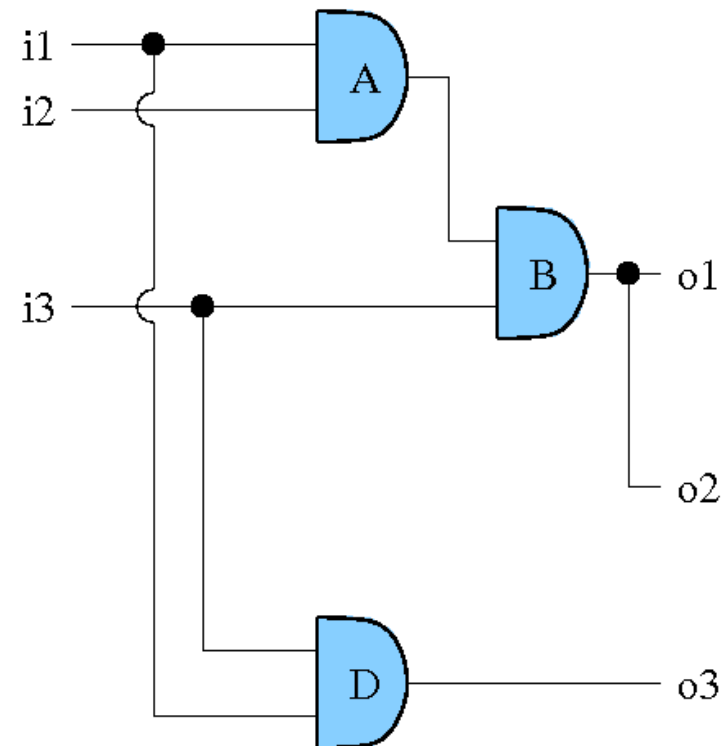
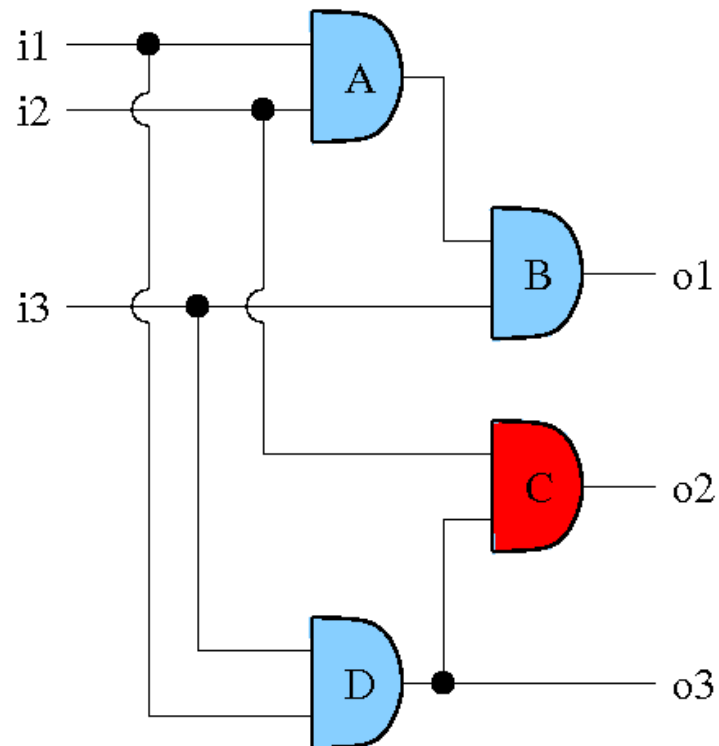
- Often *not* a substantial cause of verification complexity
 - Most of the design is insensitive to this metric
- Verification algorithms can often treat such arrays more abstractly with *memory consistency constraints*
 - Efficient Memory Model, BAT ICCAD'07, Bjesse FMCAD'08
- Though with larger caches and *more elaborate associativity schemes* comes increased complexity
 - Sometimes the logic *adjacent to* memory array becomes more complex

Moore's Heirlooms: *Circuit and Device Cleverness*

- Countless tricks behind increasing MIPS and computing power
 - Some of these are HUGE causes of verification complexity
- First consider techniques for circuit speed
 - Integration, interconnect speedup, miniaturization, datapath widening all eliminate speed barriers
 - Natural push to *speed up* core processing circuitry
 - How is this achieved?

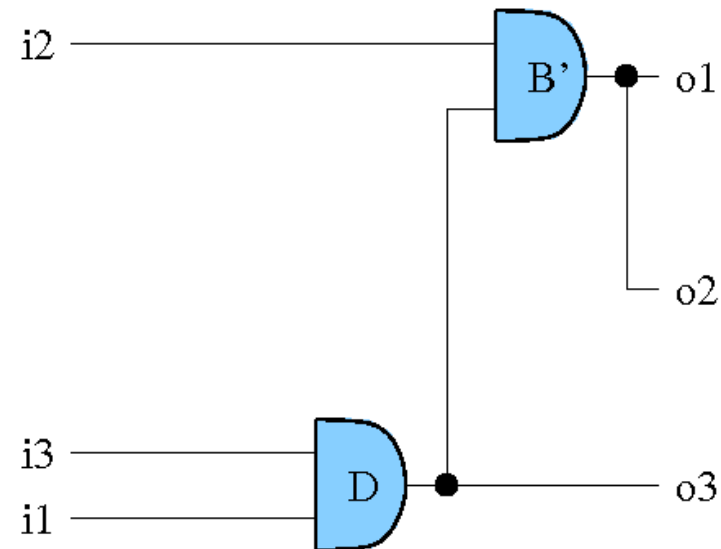
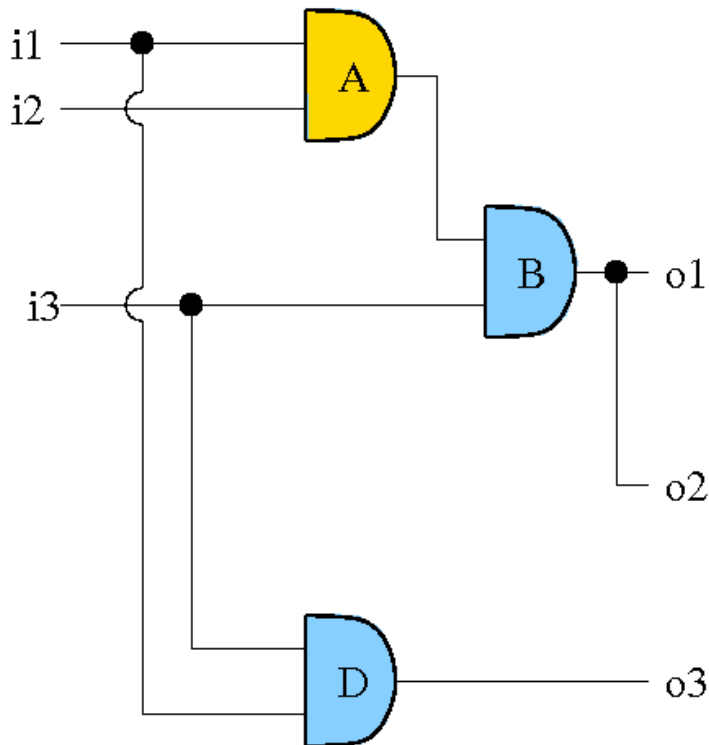
Moore's Heirlooms: *Circuit Speed*

- Decades of synthesis research to reduce logic area, delay, ...
 - E.g., simple redundancy removal and rewriting



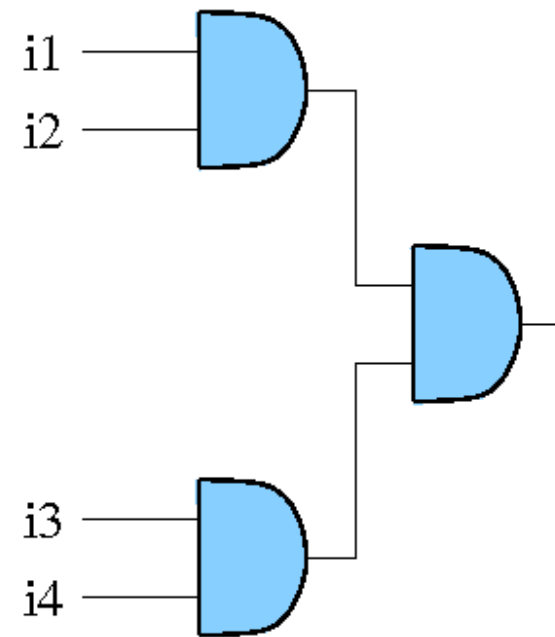
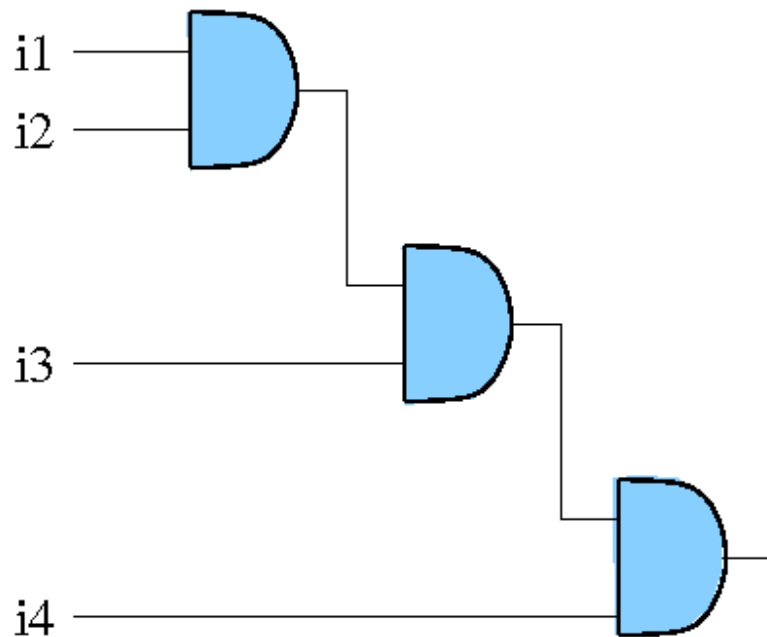
Moore's Heirlooms: *Circuit Speed*

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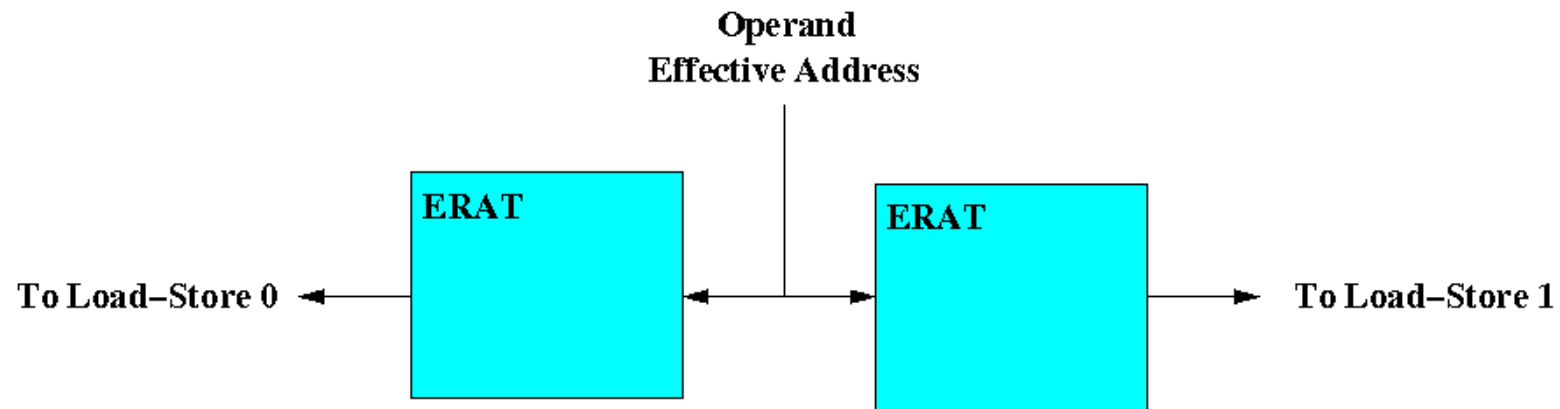
Moore's Heirlooms: *Circuit Speed*

- Decades of synthesis research to reduce logic area, delay, ...
 - E.g., simple redundancy removal and rewriting
 - Max clock speed \sim max combinational delay between state elements



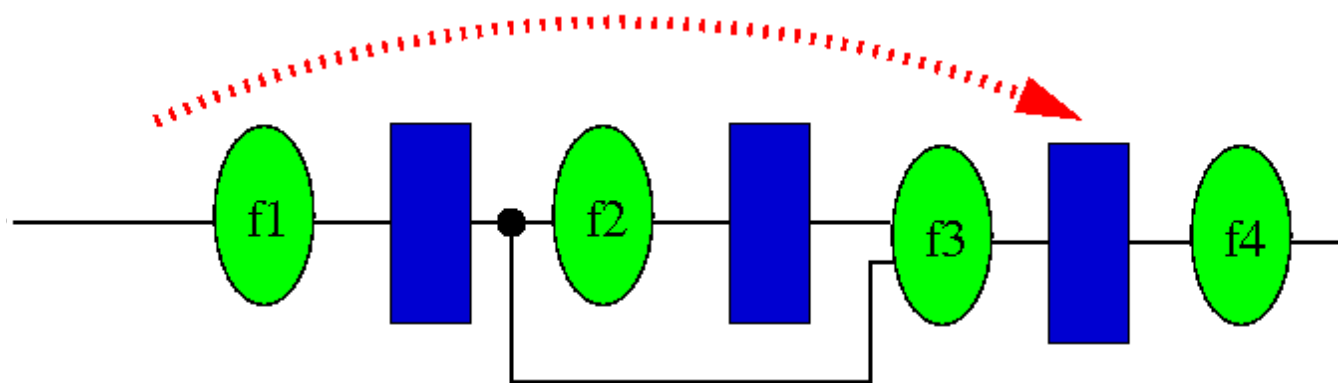
Moore's Heirlooms: *Circuit Speed*

- Decades of synthesis research to reduce logic area, delay, ...
 - Redundancy *introduction* may also speed up circuit



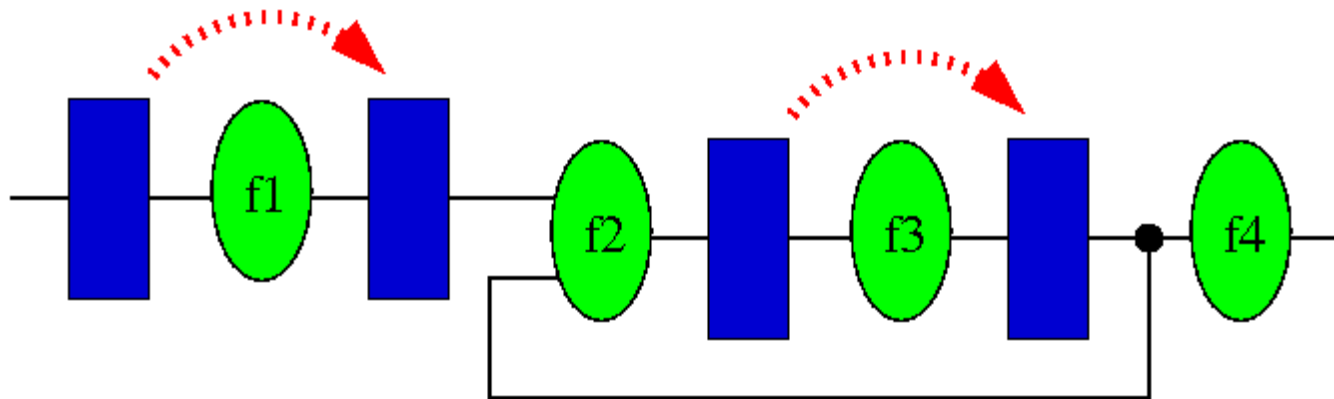
Moore's Heirlooms: *Circuit Speed*

- Retiming may be used to enable higher circuit speed
 - Max clock speed \sim max combinational delay between state elements



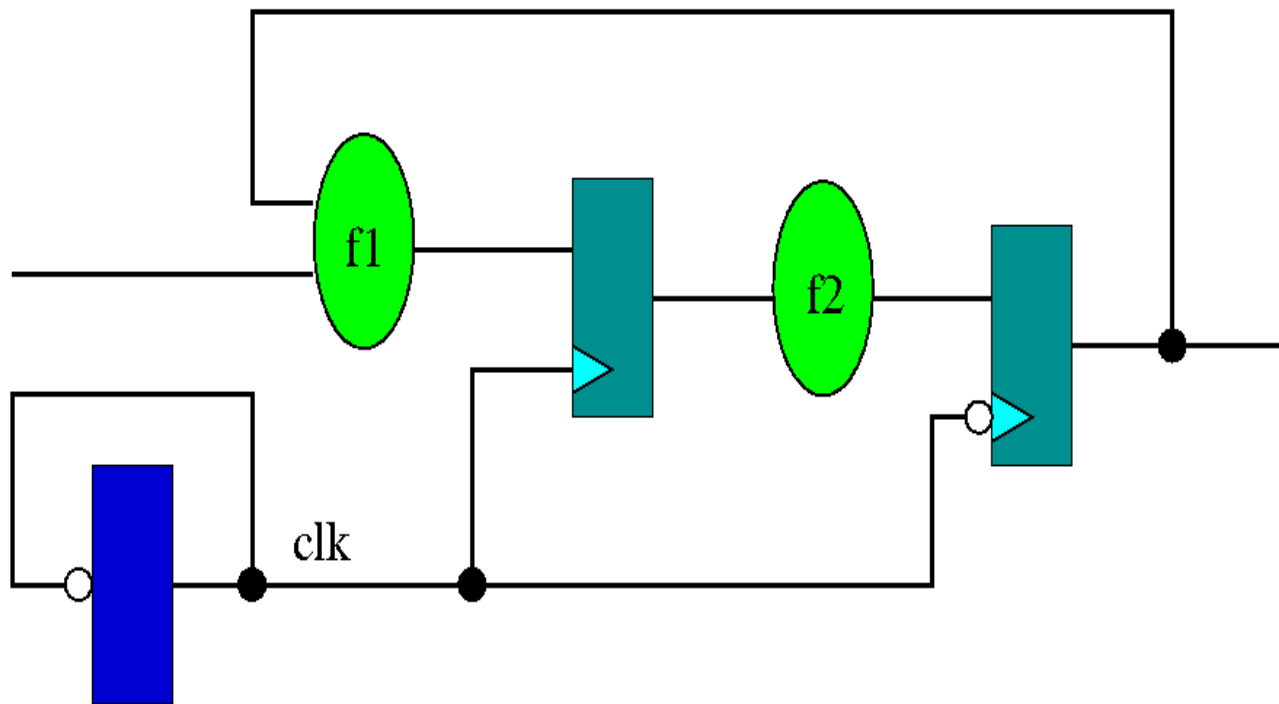
Moore's Heirlooms: *Circuit Speed*

- Pipelining (+retiming) may be used to break lengthy computations



Moore's Heirlooms: *Circuit Speed*

- Multi-phase latching/clocking better distributes propagation delays
 - Convert edge-sensitive to level-sensitive state elements; retime



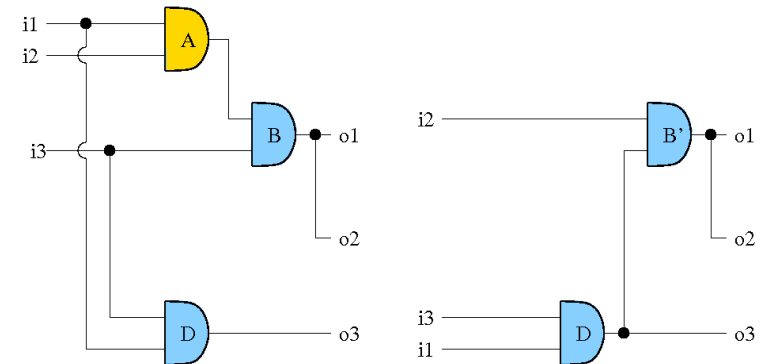
Moore's Heirlooms: *Circuit Speed*

- Do these techniques *necessarily* hurt verification complexity?
 - Some may, others not
- Luckily, many are *reversible* by verification-helping transforms

Moore's Heirlooms: *Circuit Speed*

■ Logic minimization techniques often *help* verification

- Though *word-level techniques* may suffer
- Optimizations often done at *bit-level*

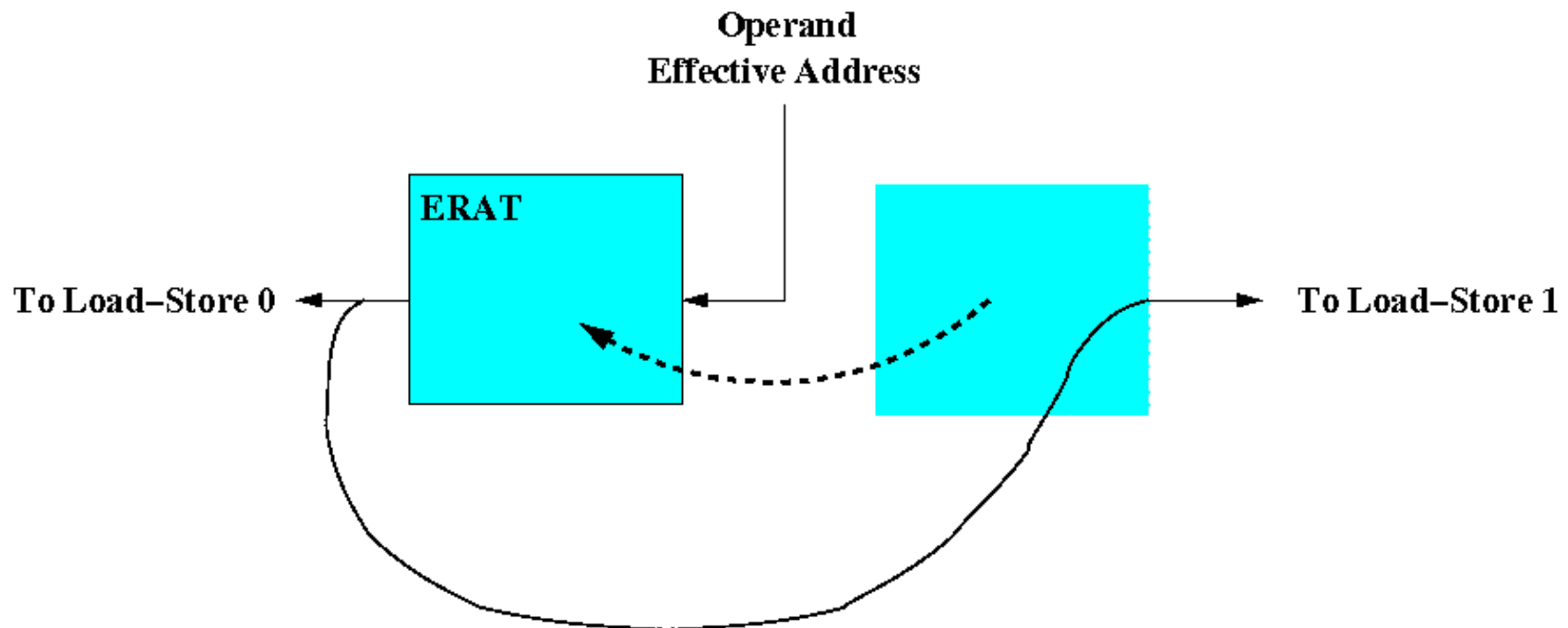


■ Indeed, often desirable to *leverage* such minimization techniques explicitly *to enhance verification*

- Minimize suboptimally synthesized designs for enhanced verification
- Exploit optimization potential created by testbench
- A powerful synergy between synthesis and verification

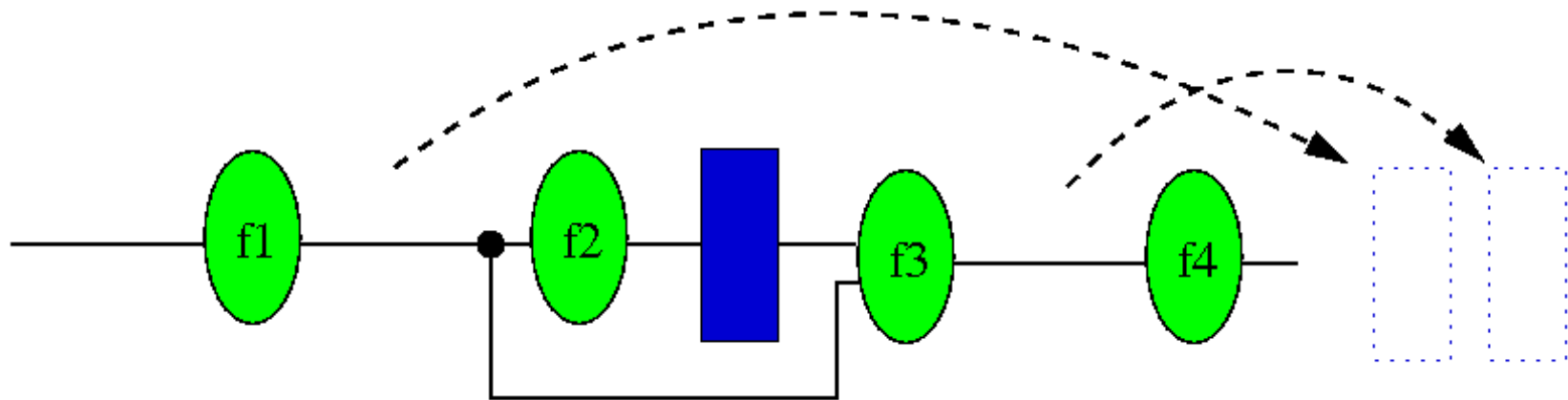
Moore's Heirlooms: *Circuit Speed*

- Redundancy *introduction* often hurts verification
 - More state elements with high correlation, more logic to reason about
 - Naive BDD, SAT algorithms may become highly inefficient
 - Induction becomes less effective with correlated state element bloat
- May be reversed through redundancy *removal*



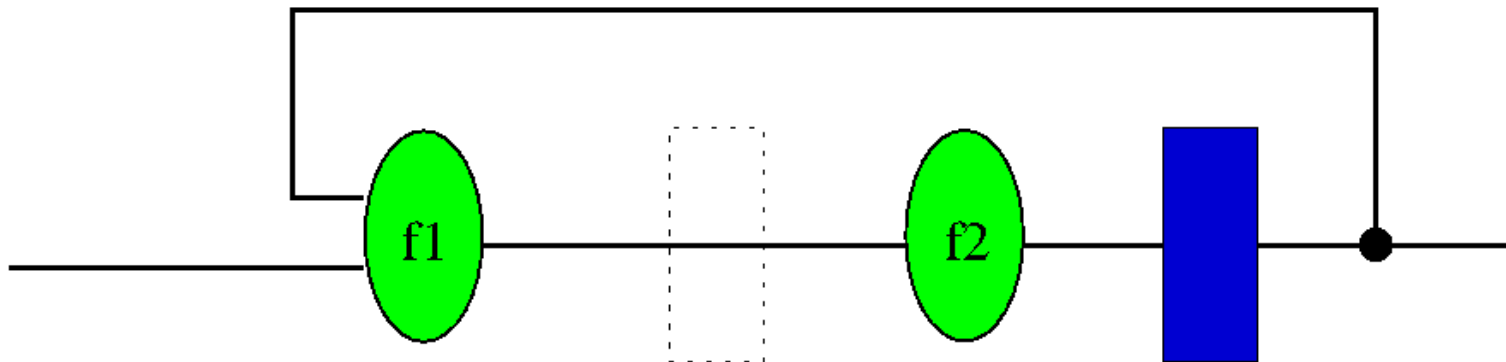
Moore's Heirlooms: *Circuit Speed*

- *Min-period* retiming often hurts verification
 - Often increases state element count, correlation
 - Similar complexities as redundancy introduction
- May be reversed through *min-area* retiming



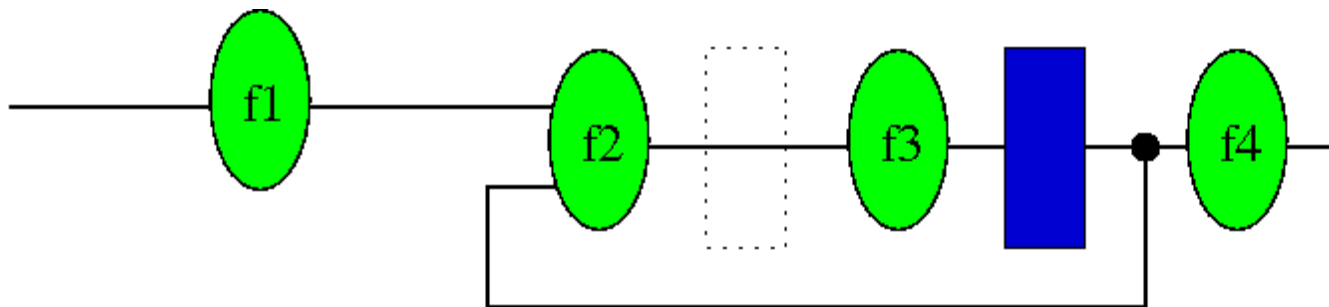
Moore's Heirlooms: *Circuit Speed*

- Multi-phase latching often hurts verification
 - Increase in state element count, correlation; increase in *diameter*
- May be reversed through *phase abstraction*
 - Unfold next-state functions modulo 2



Moore's Heirlooms: *Circuit Speed*

- Pipelining often hurts verification
 - Increase in state element count; increase in diameter
- May be reversed through peripheral retiming
 - And state-folding abstraction (like phase abstraction)



Moore's Heirlooms: *Circuit Speed*

- Overall message: correct-by-construction synthesis transforms can often be reversed to avoid significant verification penalty
 - Will be revisited later in context of “equivalence checking”
- Such transforms often *help* verification in themselves
 - ABC (UC Berkeley) won the 2008 HWMCC largely due to such transforms
 - SixthSense (IBM) uses such a transformation-based verification paradigm

Moore's Heirlooms: *Circuit Speed*

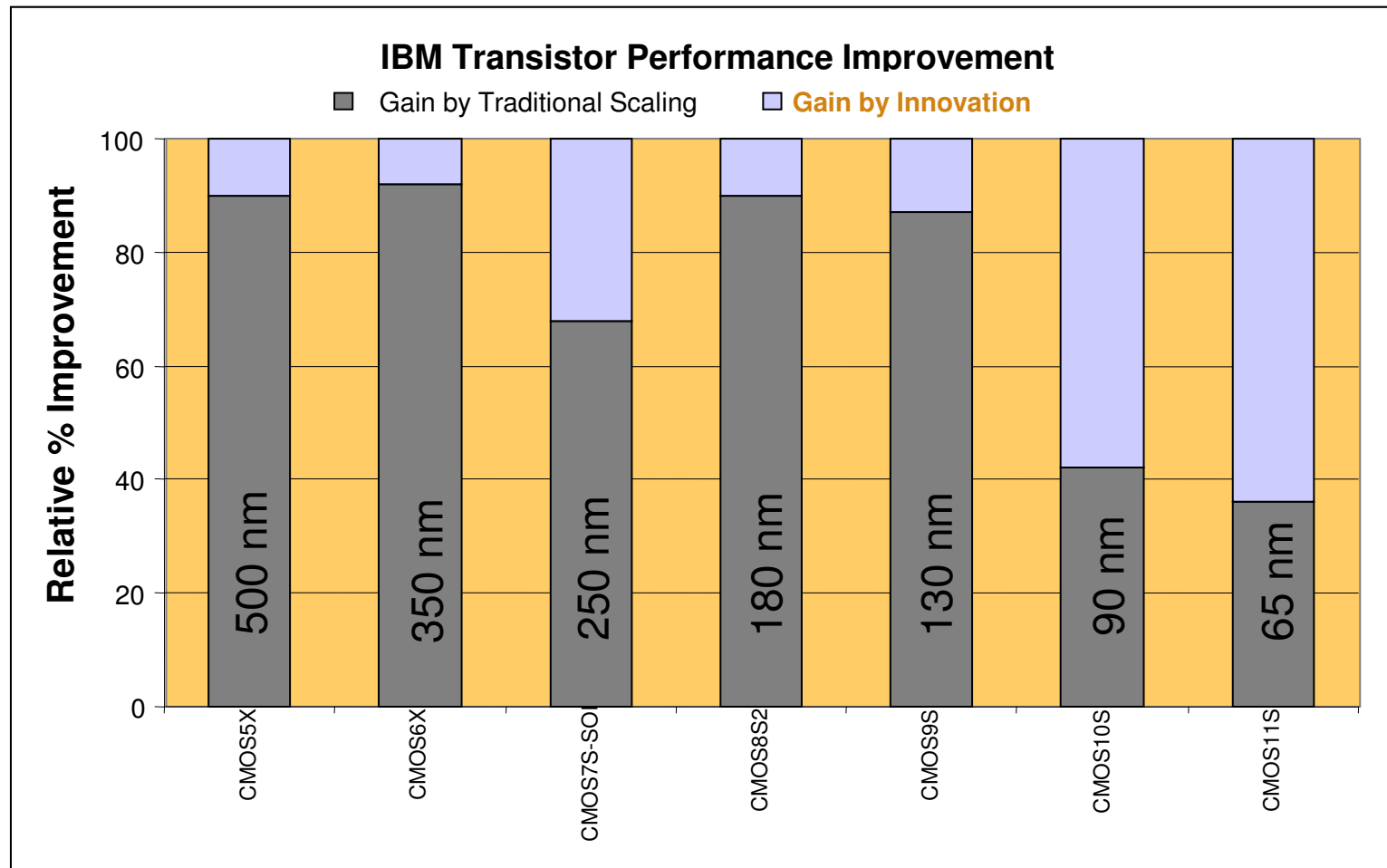
- Can *all* circuit speedups be reversed to avoid verif complexity?
 - Unfortunately (and obviously), no!
- E.g., pipelining is often accompanied by *bypass / stall* circuitry
 - Only “correct” under specific assumptions / handshaking with adjacent logic
- Entails significant design complexity
 - Cannot generically be reversed without expensive *global* reasoning

Moore's Heirlooms: *Design Cleverness*

- Many techniques critical to preserving Moore's Law momentum
 - Yet very complex for verification
 - Difficult to automate or generically “reverse”
- Superscalar + out-of-order execution
- Prefetching
- Speculative execution
- Reconfigurable hardware
- *Holistic design*
 - The simultaneous optimization of:
 - materials, circuits, cores, chips, system architecture, software, ...

Moore's Heirlooms: *Design Cleverness*

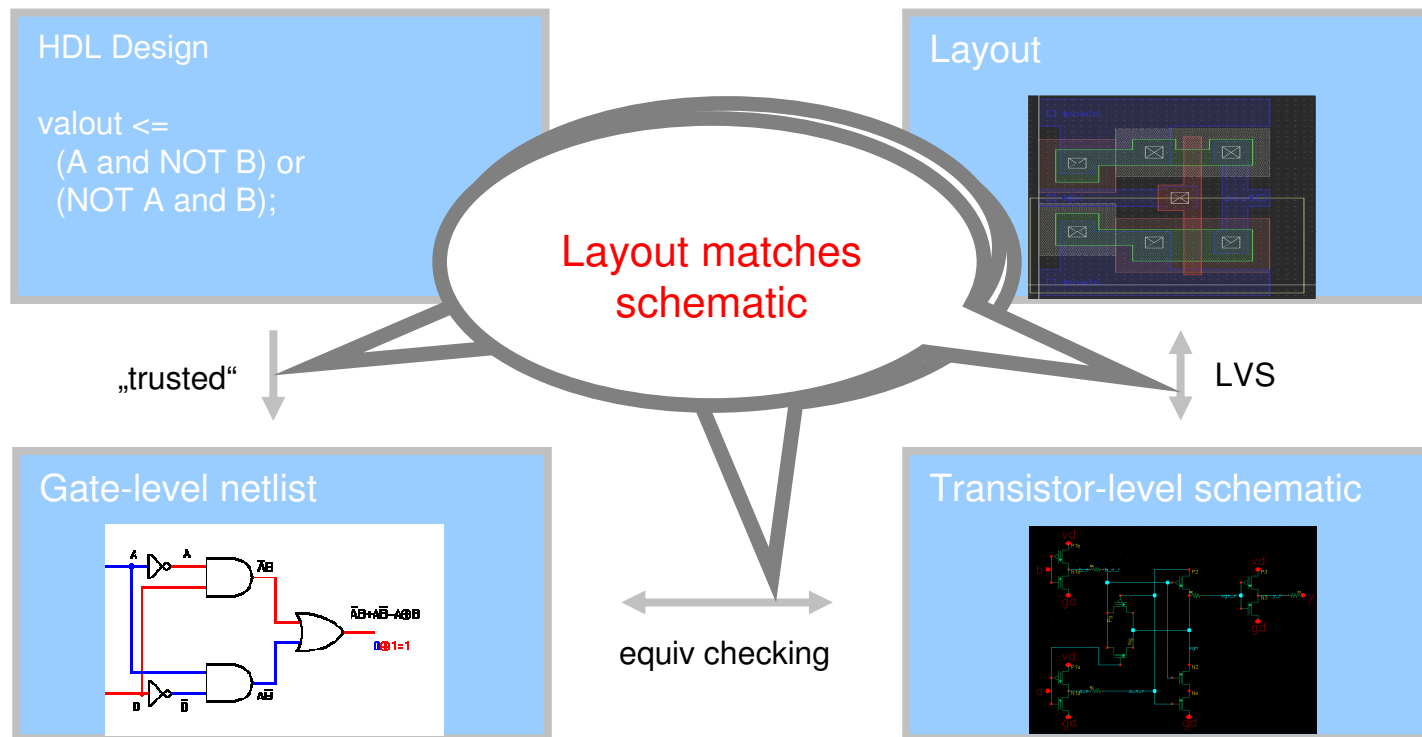
Design innovation has become key to performance gains



Must the Verif Engineer suffer Moore's Prodigy?

- Discussed circuit / design tricks which keep Moore's Law moving
 - Impact on verification
 - Ways to *attempt to* cope
- For HDL verification, why are (some of) these even pertinent??
 - Does the *HDL* need to reflect *circuit* optimizations?
- Let us briefly discuss path from HDL to fabrication

From HDL to Fabrication

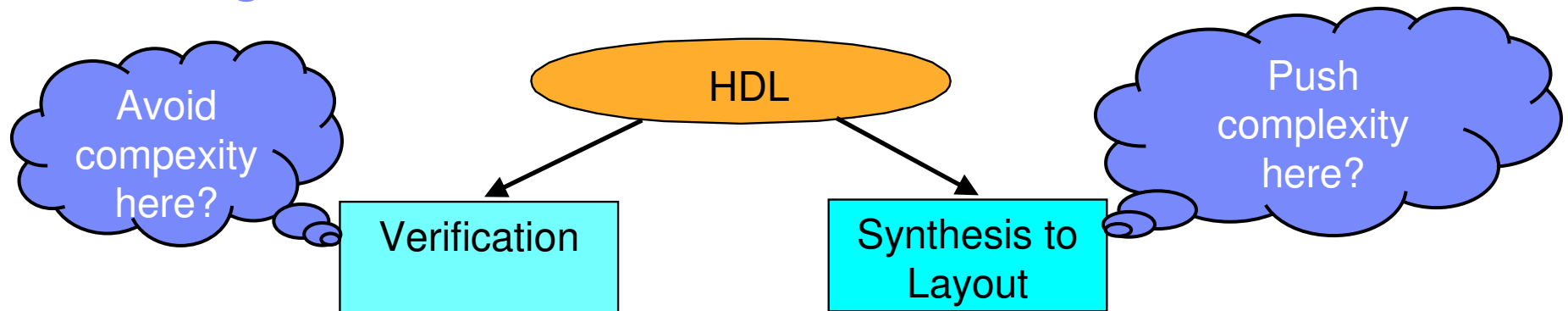


Must the Verif Engineer suffer Moore's Prodigy?

- Logic synthesis includes numerous optimizations to improve schematic quality
 - Technology-independent optimizations
 - Technology-dependent optimizations
 - Technology mapping

- Why not use a *higher-level HDL* as basis of verification?
 - Leave the ugly circuit optimizations to synthesis and equiv checking!!!

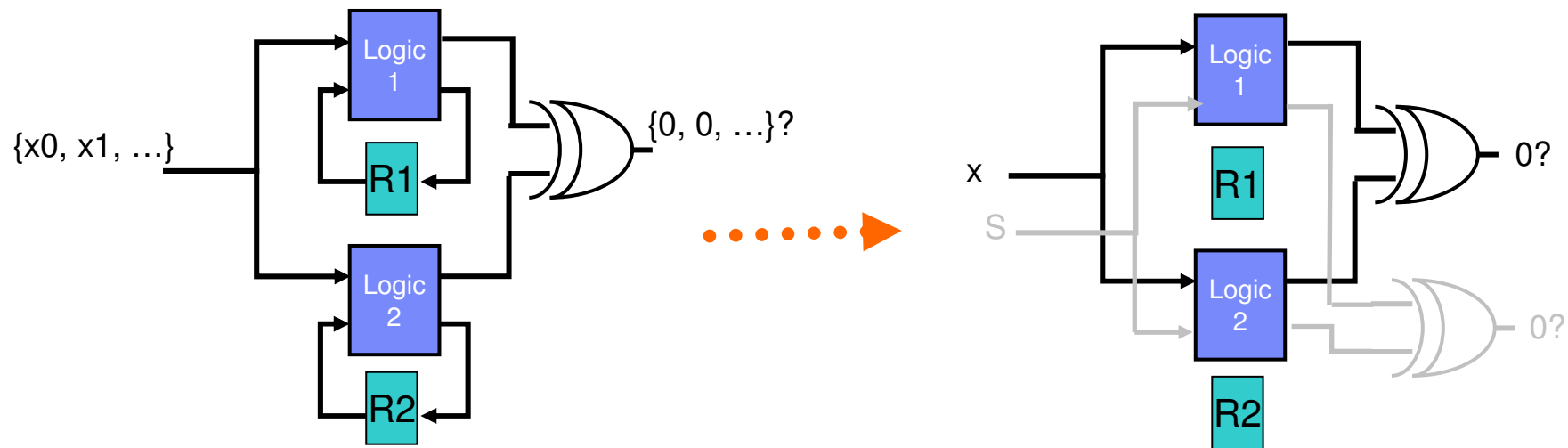
Use Higher-level HDL Basis for Verification?



- + Easier to design
- + Easier to verify
 - Simpler; amenable to word-level techniques
- + Likely to be less buggy
 - Bugs / lines of HDL roughly constant across design generations
- Any drawbacks??

Combinational Equivalence Checking (CEC)

- Very well-established technology
 - ~Every chip fabricated today leverages this technology



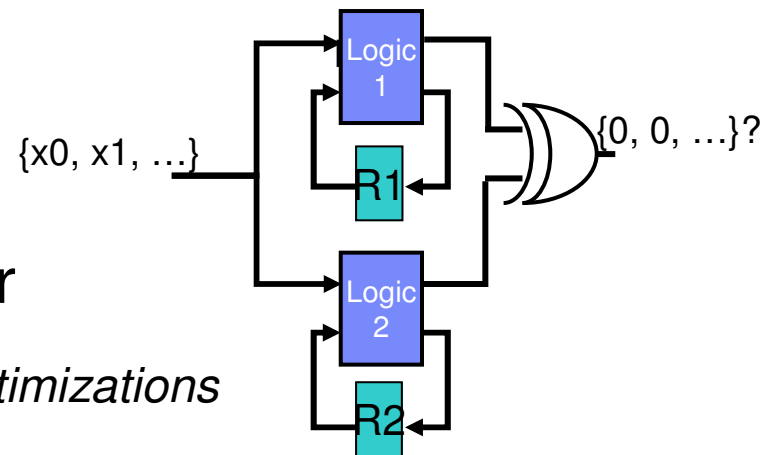
- Requires 1:1 state elt mapping between Pre / Post Synthesis models
 - Does not truly analyze sequential behavior of the design
 - Validate next-state functions, outputs with respect to cutpoints

Combinational Equivalence Checking (CEC)

- Combinational optimizations easily pushed into synthesis flow
 - No need to hand-tune HDL for combinational optimization
 - (As long as synthesis is powerful enough)
- Though not applicable for sequential optimizations
 - Retiming, pipelining, multi-phase latching, clock gating, sequential redundancy elimination/introduction, ...
 - Many CEC tools have (very) limited support for some of these

Sequential Equivalence Checking (SEC)

- Is there a technology for equiv checking *sequential* optimizations?
 - Of course! SEC has been proposed 16 years ago
 - Only beginning to see industrial application



- Analyzes true sequential I/O behavior
 - Supports *arbitrary functionality-preserving optimizations*
- With this generality comes computational complexity

Sequential Equivalence Checking (SEC)

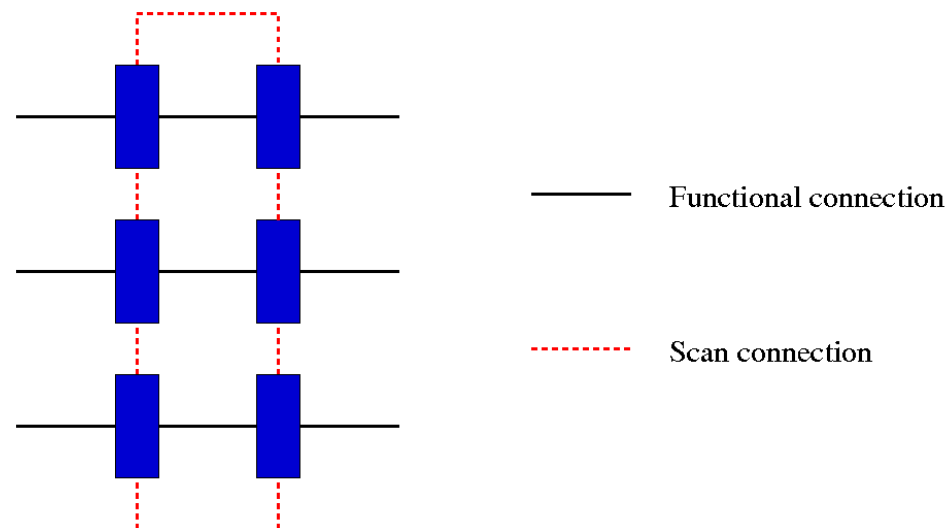
- Can we push sequential optimizations into synthesis flow?
- Due to SEC, *in theory* the answer is “yes”
- If synthesis is *powerful enough*, in practice the answer is “yes”
 - (Or is it?)

Sequential Equivalence Checking (SEC)

- Can we push sequential optimizations into synthesis flow?
 - *Optimal synthesis is as computationally expensive as verification*
 - Holistic design: synthesis alone is not yet *nearly* sophisticated enough
 - Many global optimizations require pseudo-manual “leap of faith” to perform at component level, clever methodology to validate
 - SEC is also computationally expensive!
 - Sequential reasoning over *two* designs
 - “Synthesis history”-aware approaches are promising to flatten this concern
 - Brayton’s FMCAD’07 keynote, Brayton/Mishchenko FMCAD’08
- *Post-silicon analysis is an open challenge*

Post-Silicon Analysis

- Various *pervasive logic* in chip aside from core functionality
- E.g., scan chains enable reading the contents of a chip



- Critical to debug silicon failures
 - Also used for purposes such as initialization, self-test, ...

Post-Silicon Analysis

- If synthesis moves the state elements, what exactly will be scanned?
 - Sequential optimizations *must* be done *before* scan insertion (else ineffective)
 - Need to remap scanned values to state elements familiar to the designer
- Can mapping be done through sim + SAT using “synthesis history”?
 - May entail additional constraints on synthesis; de Paula FMCAD’08, Hunt
- What about trace/debug buses, which monitor the chip *real-time*?
 - Cannot algorithmically post-process that trace!
- What about Engineering Change Orders?
 - Last-minute changes made directly to post-synthesis design

Post-Silicon Analysis

- If we can reconcile pervasive design w.r.t. sequential synthesis
- We still need to *verify* the pervasive logic
 - 1) Verify that primary functionality is unaltered through this injection
 - Easy SEC-style task, constraining to functional analysis
 - Though now synthesis *does* conditionally alter behavior!
 - 2) Need to verify that pervasive logic works properly!
 - Often more critical than main functionality; chip rendered useless otherwise
- What if not correct (or inadequate): how debugged?
 - Does the designer need to learn post-synthesis design ??
 - Somewhat defeats the purpose !!!

Coping with Moore's Heirlooms

- How can verification cope with increased design complexity?
 - *A good methodology is often key overcoming technology shortcomings*

1) Abstraction?

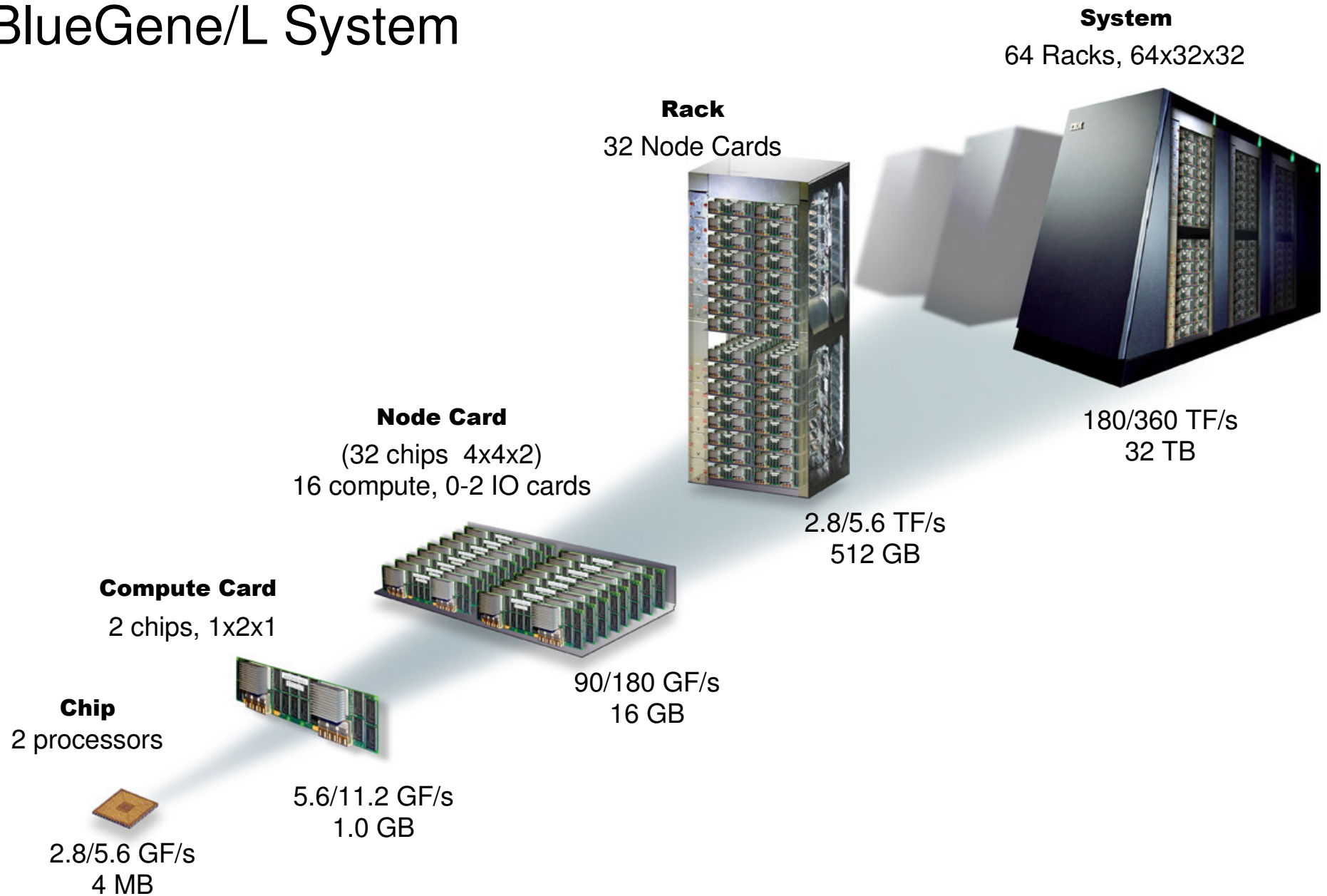
- Maybe, though intrinsic complexity need to be modeled+verified *somewhere*
 - *Leaky abstractions* have their limits!



Coping with Moore's Heirlooms

- *Leaky abstractions* have their limits
- Push for *holistic design* makes these limits fundamental
- May seem crazy, though a cost-effective path to high-end systems
 - Suppressing design detail too deeply risks catastrophe
 - Failure to meet timing / performance / power / reliability goals
 - *Lost time to market*
 - **Product failure**

BlueGene/L System



Coping with Moore's Heirlooms

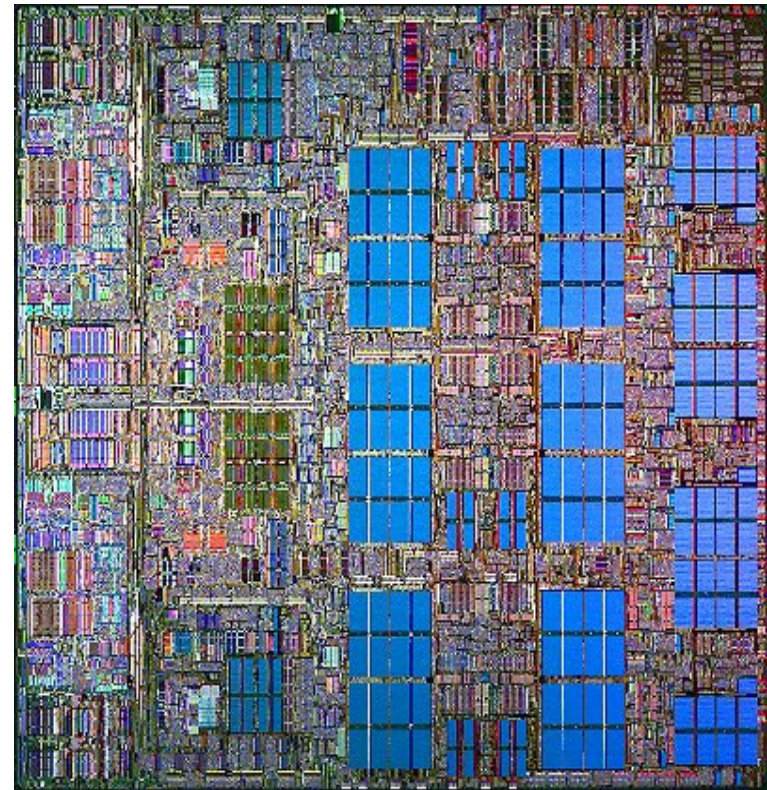
- *Leaky abstractions* have their limits
- Though abstraction clearly has a substantial role in verification
 - May enable a *hierarchical verification process*
 - *Protocol* verification → *HDL* proof obligations
 - *Unoptimized design* verification → *optimized design* proof obligations
- Sometimes requires *decomposition* of verification process
 - E.g. functionality vs pervasive vs timing vs power
 - Much harder to decompose *design* process in this manner, unfortunately
- Manually maintaining abstractions is *expensive*
 - Higher-level HDL, automated optimizations are “free abstractions”

POWER5 Chip

- Dual pSeries CPU
- SMT core (2 virtual procs/core)
- 64 bit PowerPC
- 276 million transistors
- 8-way superscalar
- Split L1 Cache (64k I & 32k D) per core
- 1.92MB shared L2 Cache >2.0 GHz
- Size: 389 sq mm
- 2313 signal I/Os
- **>>1,000,000 Lines HDL**
 - **How big would its abstractions be?**

POWER architecture:

- Symmetric multithreading
- Out-of-order dispatch and execution
- Various address translation modes
- Virtualization support
- Weakly ordered memory coherency

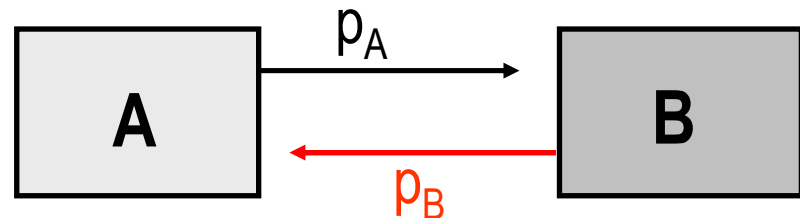


Coping with Moore's Heirlooms

- How can verification attempt to cope with increased design complexity?

2) Compositional reasoning

- Isolate complex logic as precisely as possible for dedicated verification
- When verifying adjacent logic (or higher-level model), abstractly model these complex components



- *Assume-guarantee reasoning* is sound and complete
 - Substantial work in automating this process
 - Still largely manual for realistic designs

Coping with Moore's Heirlooms

- Compositional verification is often manually intensive
 - Unless design is well-partitioned with well-defined interfaces
- For highest performance, *global optimization* is often performed
 - Subtle signal / timing relations are exploited between blocks
 - Intricate functionality decentralized
- Optimization to the point that designs *almost don't* work properly
- Overall: Even clever methodologies become cost-ineffective to cope with technology shortcomings
 - In cases need to resort to incomplete techniques; risk verification gaps

Example: *Power Saving Logic*

- Higher speed, increased density → need for power reduction!
- Karen already spoke about power-gating logic
 - Reduce power consumption by disabling voltage to *inactive* components
- Methodologically can cope with this complexity
 - 1) (Compositionally) verify that it does not meaningfully alter functionality
 - 2) Disable power-gating logic for ease of functional verification
 - A simple “abstraction”
 - 3) Also need to verify that it works as intended!
- *Decompose and abstract*

Example: *Power Saving Logic*

- Asynchronous issues arise due to clock domain decoupling
 - Even if running at identical frequencies
- *Metastability* may propagate if not properly guarded against
- Communication channels may suffer sequential delays
- Synchronization logic used to protect design from these dangers
 - Another set of verif tasks to ensure that this logic works properly
- “Abstract design” attained by not injecting metastability, delays
 - Though still includes the complexity of the synchronization logic
 - Manual abstraction / higher-level design may be needed



Conclusion

- Does Moore's Law hurt the verification community?
 - Yes!
 - Though not always!
- *Many open problems in design + verification*
 - *HW verification is not a solved problem*
- *Old open problems:*
 - Improve bit-level verification algorithms !
 - Improve bit-level synthesis algorithms !
 - Improve equivalence checking techniques !



Conclusion

■ *Newer* open problems

- Improve higher-level verification algorithms !!
- Improve high-level synthesis techniques !!
- Improve higher-level equivalence checking techniques !!

■ *Newest* open problems

- Improve design+verification methodologies w/ focus on holistic design !!!
- Address post-Si issues through sequential synthesis !!!

■ “Someday Moore’s Law will work *for*, not *against*, the verification community” Allen Emerson

- Requires substantial innovation! *Help achieve this goal !!!!*